Data Integration and Large Scale Analysis

Slides credit: Matthias Boehm - Shafaq Siddiqi

10- Distributed Data-Parallel Computation



Lucas Iacono. PhD. - 2024



Part B

Large-Scale Data Management & Analysis

- LU3. Cloud Computing
 - Cloud Computing Fundamentals
 [Nov 29]
 - Cloud Resource Management and Scheduling [Dec 06]
 - Distributed Data Storage[Dec 13]

Part B

Large-Scale Data Management & Analysis

- LU4. Large-Scale Data Analysis
 - Distributed, Data-ParallelComputation [Dec 20]
 - Distributed Stream Processing
 [Jan 10]
 - Distributed Machine Learning Systems [Jan 17]

Agenda

- Announcements
- Data-Parallel Collection & Processing
- Data-Parallel DataFrame Operations

Announcements

Announcements

Course Evaluation and Exam

- Course evaluation: 20/02/2025
- Exam date: Feb 07, 3:00pm (90 min written exam)
- Oral Exam for Erasmus Students
 - Schedule available in TeachCenter (23/12/2024)

Recap: Distributed Collections	Кеу	Value
Logical multi-set (bag) of key-value pairs (unsorted collection)	13:00:01	12.1
Different physical representations key-value pairs can be	14:00:05	16.0
stored in various ways (e.g., database, across files, or in memory).	13:00:03	12.5
Easy Distribution via Horizontal Partitioning. Data divided	13:00:05	13.0
into "chunks" (shards or partitions) based on the keys. Each chunk stored on a different machine (easier to handle	14:00:04	15.7
large-scale data).	14:00:06	16.3
How collections are created: from single file with data or a	13:00:00	12.1

folder of files (even if they're messy and unsorted).

Recap: Files and Objects

File: large and continuous block of data saved in a specific format (CSV, Binary, etc.).

Object: like a file, but binary and it comes with metadata (Images on S3)

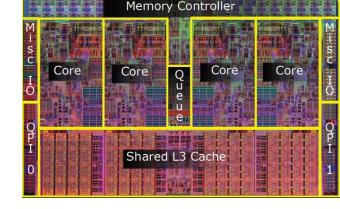
Recap: Object Storage

1. Object Storage (e.g. AWS S3):

- a. Data stored as objects (data, metadata, and UID).
- b. Ideal for storing unstructured data like media files, backups, or large datasets.
- c. Objects of a limited size (e.g., 5TB in AWS S3).

Nehalem Architecture

- Integrated Memory Controller: Integrated in chip, -- latency and ++ memory performance.
- Support for DDR3 Memory: Higher memory bandwidth (compared to DDR2).
- QuickPath Interconnect (QPI): High-speed, point-to-point connection (no Front-Side Bus).
- Enhanced Hyper-Threading: Each core supports two threads (+++ performance)
- Multi-Core Scalability: 2 to 8 cores per processor (2 threads / core)
- Improved Cache Design: Dedicated L1 and L2 cache p/core shared L3 cache





Michael E. Thomadakis: The Architecture of the Nehalem Processor and NehalemEP SMP Platforms, Report, 2010

Nehalem Architecture

- Energy Efficiency: Turbo Boost for dynamic clock speed adjustments.
- Advanced Manufacturing Process: Higher transistor density and better efficiency.
- Integrated Graphics (in later models): Some models included integrated GPUs.
- Foundation for Modern Architectures: Established the groundwork for subsequent Intel architectures like Sandy Bridge and Skylake.

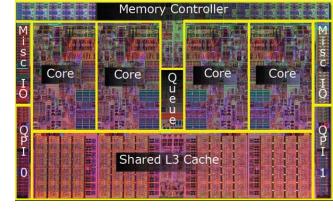


Memory Controller

Michael E. Thomadakis: The Architecture of the Nehalem Processor and NehalemEP SMP Platforms, Report, 2010

Nehalem Architecture

- Pipeline
 - Frontend:
 - Instruction Fetch
 - Pre-Decode
 - Decode CISC 2 uOps (ADD [eax], 5)
 - Load the value from memory.
 - Add 5 to the loaded value.
 - Store the result back to memory.
 - Backend:
 - Rename/Allocate
 - Scheduler
 - Execute

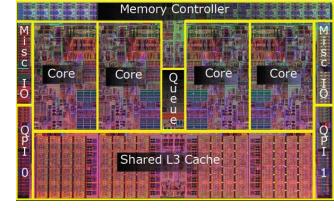




Michael E. Thomadakis: The Architecture of the Nehalem Processor and NehalemEP SMP Platforms, Report, 2010

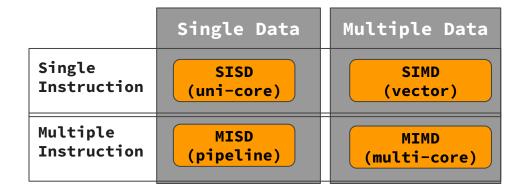
Nehalem Architecture

- Out-of-Order
 - Instructions are not necessarily executed in the order they appear in the program
- Execution Engine: 4 Inst x Cycle (IPC=4)
- 128-bits Floating-point multiplication
- 128-bits floating-point addition





Michael E. Thomadakis: The Architecture of the Nehalem Processor and NehalemEP SMP Platforms, Report, 2010



Flynn's Classification

Computer architectures based on how they handle instructions and data.

- SISD:
 - One task at time one data chunk (e.g. PC running a single program)
- SIMD:
 - One task at time multiple data chunks (e.g. GPUs rendering)
- MISD:
 - Multiple tasks one data chunk (e.g. fault-tolerant computers)
- MIMD:
 - Multiple tasks multiple data chunks (multi-core CPUs 1 Core -> Program)



Michael J. Flynn, Kevin W. Rudd: Parallel Architectures. ACM Comput. Surv. 28(1) 1996

Distributed, Data-Parallel Computation

- Parallel computation of function **foo()** → **single instruction**
 - A single function applied to all data items in parallel.
- Collection X of data items (key-value pairs) → multiple data
 - foo() operates on multiple pieces of data (key-value pairs).
- Data parallelism similar to SIMD but more coarse-grained notion of "instruction" and "data" → SPMD (single program, multiple data)

Y = X.map(X -> foo(x)) X = Data Items (e.g. array), .map (operation to each element in X), Y = Output



[Frederica Darema: The SPMD Model :

Present and Future. PVM/MPI 2001]

Past.

SPMD

_ ___ __

- Dynamic Work Assignment. Processes can self-schedule, ++ flexibility & efficiency.
- More General than SIMD. SPMD allows different instruction streams for different data. It can handle more complex tasks.
- Efficient Control. Performed at the application level rather than the OS level (less costly and more efficient than F&J.
- Applications:
 - MPI (Message Passing Interface)
 - PVM (Parallel Virtual Machine)
 - Grid Computing



[Frederica Darema: The SPMD Model : Past, Present and Future. PVM/MPI 2001]

Model	Key Features	Pros	Cons
BSP (Bulk Synchronous)	Global barriers enable synchronization after each phase	+++ Correctness and consistency; simple to implement	Overhead due to waiting at barriers Slow for stragglers
ASP (Asynchronous Parallel)	Processes run independently	Faster execution (no waiting)	Accuracy issues from outdated data
SSP (Stale-Synchronous Parallel)	Controlled staleness allows fastest processes to proceed within a limit	Balances efficiency and consistency; reduces waiting time compared to BSP	Small inaccuracies

Data-Parallel Collection & Processing





Brief Hadoop History

- Google's GFS + MapReduce [ODSI'04] -> Apache Hadoop (2006).
- Apache Hive (SQL), Pig (ETL), Mahout (ML), Giraph (Graph)

Hadoop Architecture / Eco System

- Management (Ambari)
- Coordination / workflows (Zookeeper, Oozie)
- Storage (HDFS)
- Resources (YARN)

Hadoop



Hadoop Ecosytem

- Apache Hive (SQL)
 - What it is:
 - Data warehouse infrastructure built on top of Hadoop.
 - Allows you to query and analyze large datasets stored in Hadoop using a SQL-like language called HiveQL.
 - Main Purpose:
 - Querying and analysis of big data using familiar SQL syntax.
 - Suitable for batch processing and data summarization.
 - Use Case:
 - Running SQL queries to analyze log data or generate business reports.

Hadoop





Hadoop Ecosytem

- Apache Pig (ETL)
 - What it is:
 - High-level platform for creating data processing programs in Hadoop.
 - **Pig language** to **simplifies the MapReduce** jobs writing process.
 - Main Purpose:
 - ETL operations. Cleaning, transforming, and preparing large datasets for analysis.
 - Use Case:
 - Processing raw web logs into structured formats for further analysis.

Hadoop





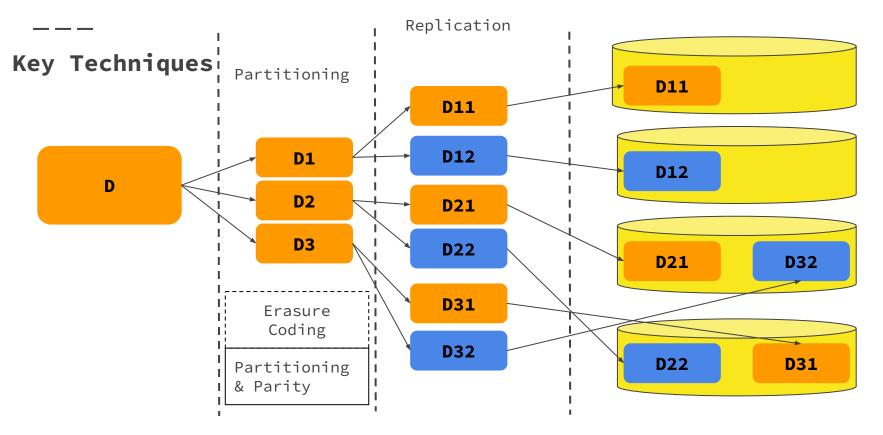


Hadoop Ecosytem

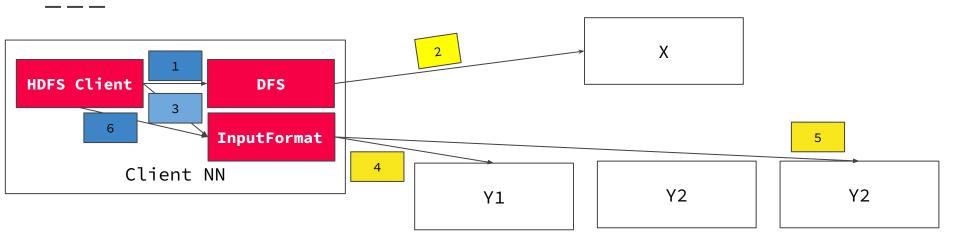
- Apache Mahout (ML)
 - What it is:
 - A library for building scalable machine learning algorithms on top of Hadoop.
 - Focused on distributed or scalable implementations of common ML algorithms.
 - Main Purpose:
 - Implementing machine learning algorithms like clustering, classification, and recommendation systems on large datasets.
 - Use Case:
 - Building a recommendation system for an e-commerce platform using collaborative filtering.



Distribution



Recap: HDFS Read

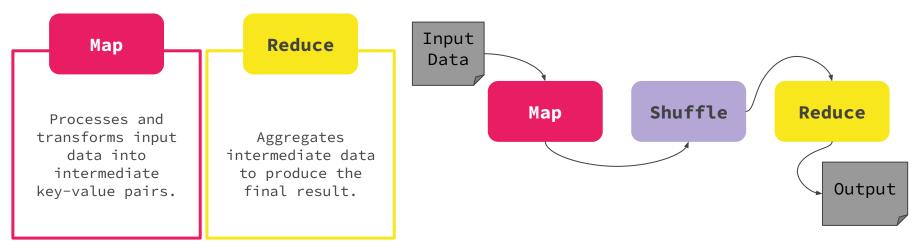


- 1. Open
- 2. Get Block Locations
- 3. Read
- 4. Read
- 5. Read
- 6. Close

MapReduce – Programming Mode

Overview

- MapReduce is a programming model for processing large datasets in parallel, distributed across multiple nodes.
- Developed by Google; popularized by Apache Hadoop.



MapReduce I

Why MapReduce?

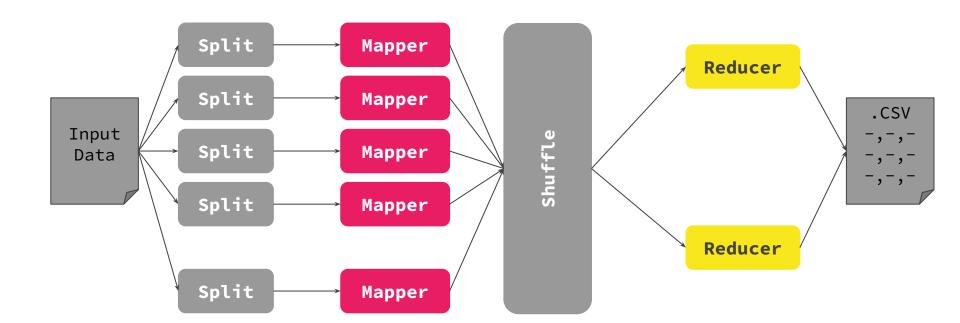
- Handles large-scale data processing efficiently.
- Works on commodity hardware.
- Built-in fault tolerance.
- Suitable for structured, semi-structured, and unstructured data.

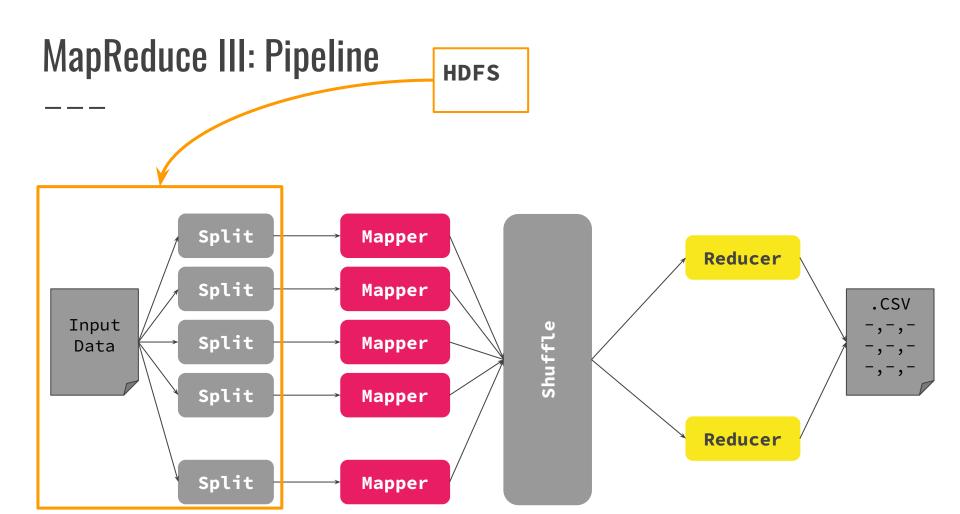
MapReduce II

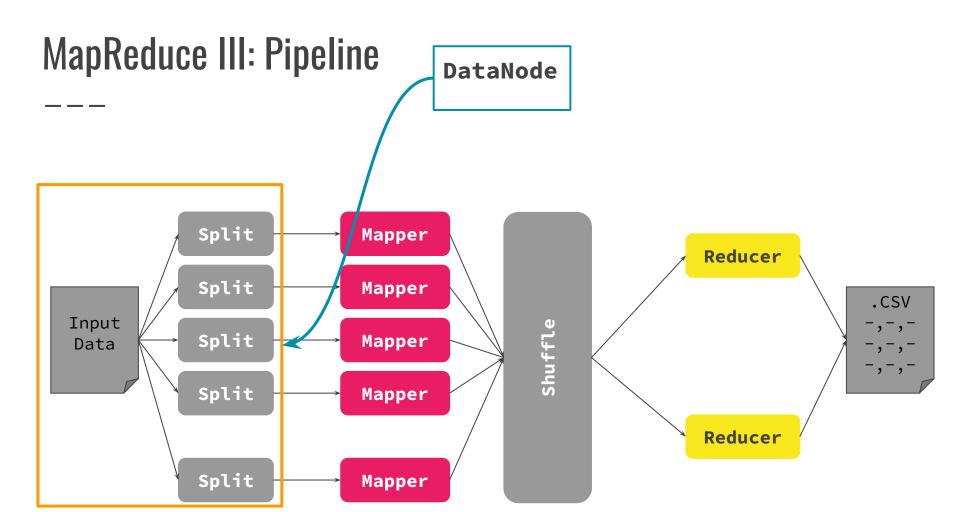
Key Concepts

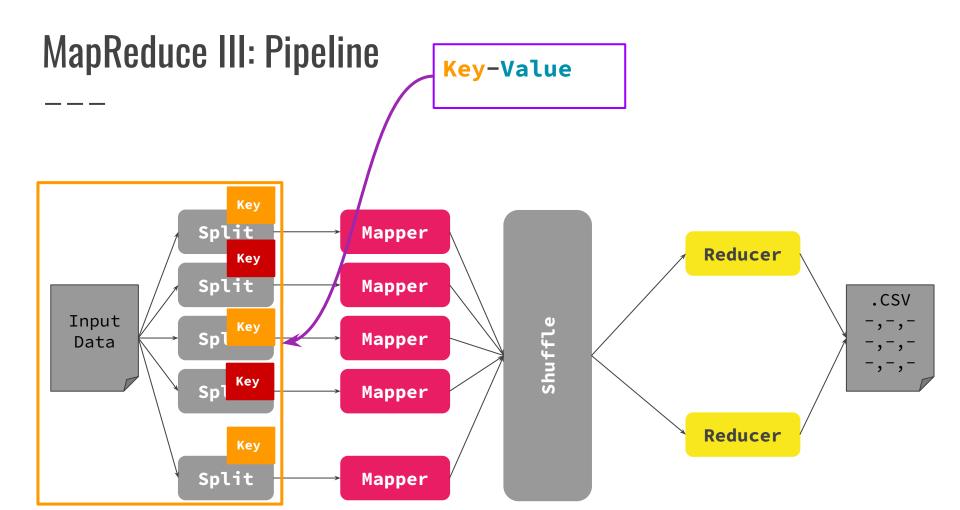
- Distributed Processing: Data is split across multiple nodes for parallel execution.
- Key-Value Pairs: Core data structure in MapReduce.

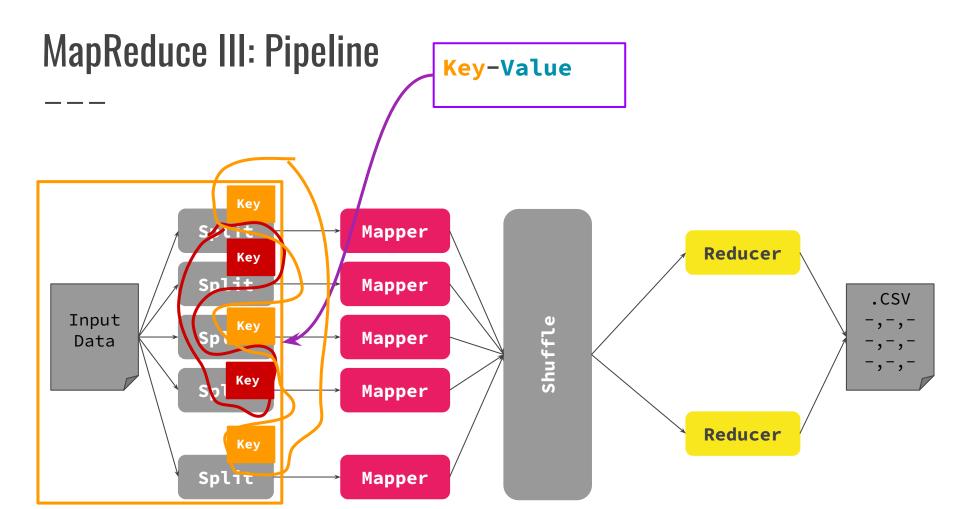
MapReduce III: Pipeline

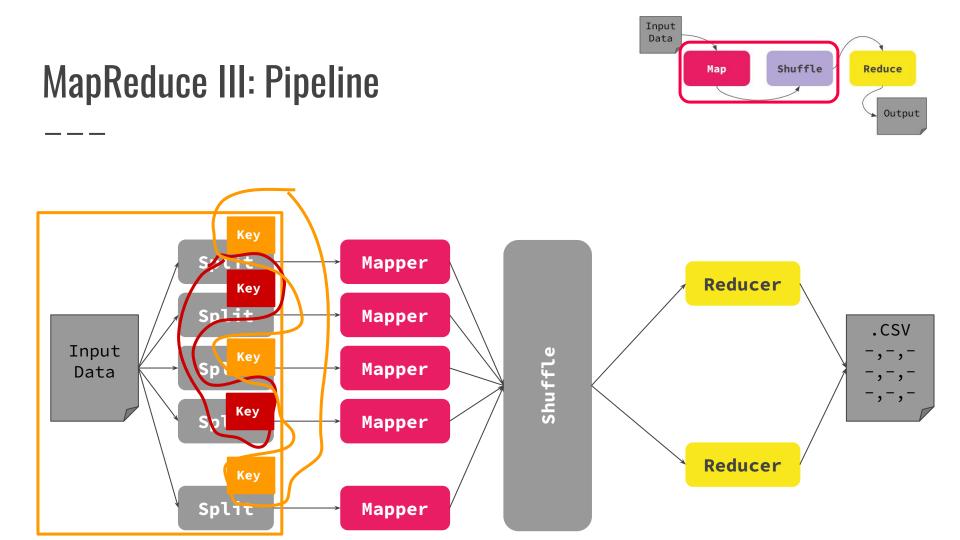




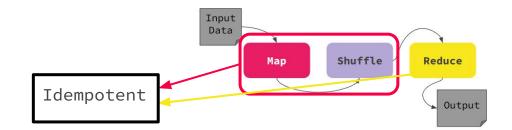


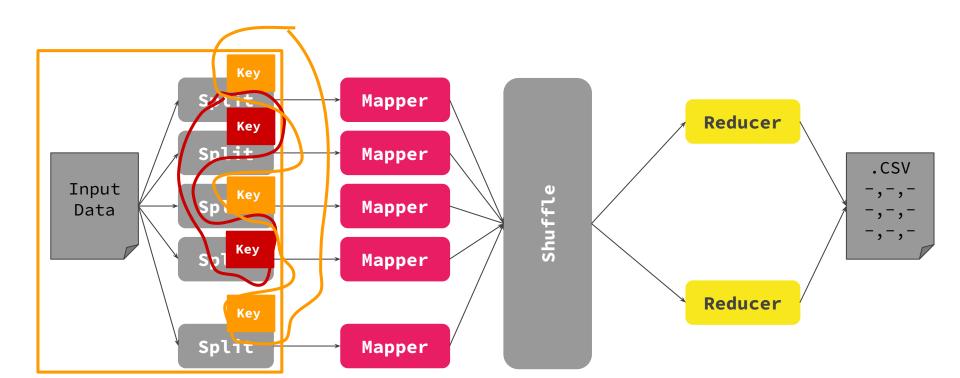


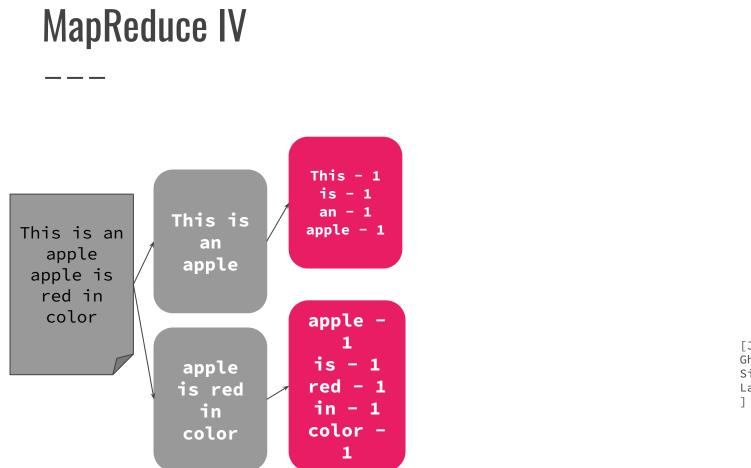




MapReduce III: Pipeline

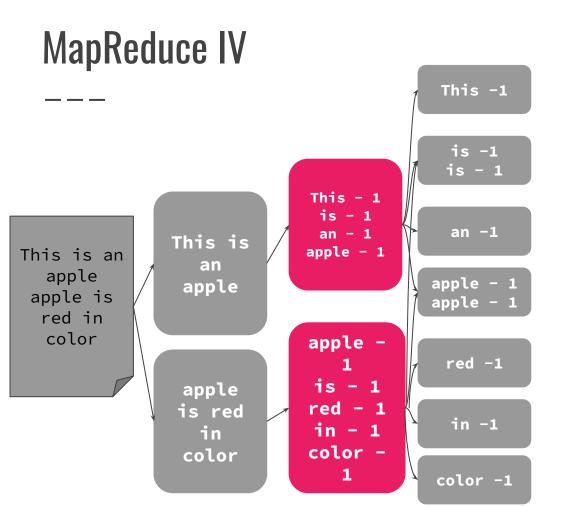


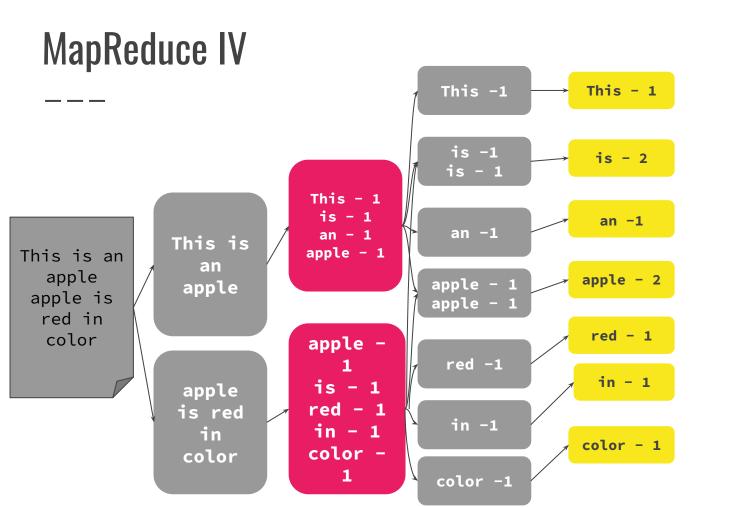




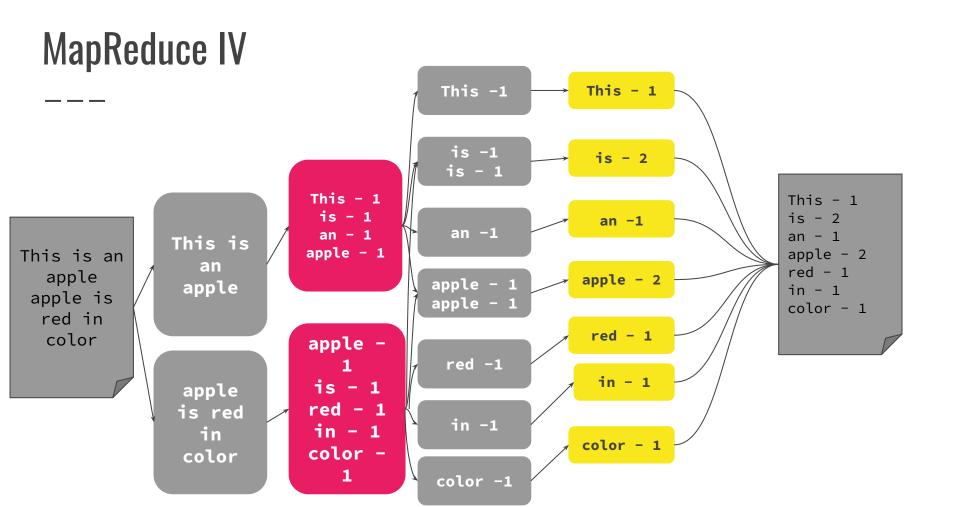


[Jeffrey Dean, Sanjay Ghemawat: MapReduce: Simplified Data Processing on Large Clusters. OSDI 2004]]





This - 1



MapReduce VI: Hands on Lab



Servers Log

- Use the MapReduce programming model to:
 - \circ Count how many times each page was accessed.
 - Identify the most popular page.
- Calculate the Average Using MapReduce
 - Given a list [4, 8, 15, 16, 23, 42], compute the average using MapReduce.

MapReduce: Summary (Pros)

_ _ _

- Large-scale processing. Large amounts of data distributed across multiple nodes in a cluster.
- Fault-tolerant. If a node fails, the system can recover and reassign tasks to other nodes.
- User Defined Functions and files. Developers can define their own custom processing logic through UDFs, and the model relies on files to store intermediate and final results.
- Flexibility. Developers can customize processing logic while the system manages distribution and fault recovery automatically.
- **Restricted functional APIs.** MapReduce relies on a limited set of functional primitives:
 - Map: Transforms input data into key-value pairs.
 - **Reduce:** Aggregates values associated with the same keys to produce results.
- **Implicit parallelism**. Developers only need to implement the Map and Reduce functions; the distribution of workload across nodes relies on the system.

MapReduce: Summary (Cons)

- **Performance:** its performance can suffer in complex workloads due to heavy reliance on I/O (writing and reading intermediate data to/from disk).
- Low-level APIs: The API is relatively basic, requiring a lot of manual effort to implement more sophisticated workflows.
- Many different systems: Specialized systems (e.g., Apache Spark, Apache Flink, or distributed database systems) have emerged as alternatives, often being more efficient and user-friendly.



Evolution to Spark (and Flink)

- Spark [HotCloud'10] + Resilent Distributed Data Sets (RDDs) [NSDI'12] → Apache Spark (2014)
- Design 1: Standing executors with in-memory storage:
 - Spark keeps **long-running worker processes** (executors) active, enabling tasks to run faster by avoiding repeated setup costs.
 - Data is stored in memory whenever possible, **minimizing disk I/O** for iterative and interactive jobs.
- Design 2: Lazy evaluation:
 - **Directed Acyclic Graph of transformations** rather than executing them immediately.
 - Actions (e.g., collect, save, count) trigger DAG's execution, allowing workflow optimization by reordering and combining operations.



- **Design 3:** Fault tolerance via RDD lineage
 - Data partition lost -> Spark can recompute using lineage graph of transformations applied to the data (reliability without heavy replication).
- Performance:
 - In-memory storage. By keeping intermediate data in memory,
 Spark significantly reduces disk I/O (faster for iterative tasks e.g machine learning).
 - Fast job scheduling. Spark's scheduler operates with low overhead, enabling tasks to be scheduled in milliseconds (~100ms), compared to Hadoop's ~10 seconds per job.



- APIs:
 - Richer functional APIs. Wide range of functional operators (e.g., map, reduce, filter, groupByKey, flatMap) compared to Hadoop -> easier to write complex workflows.
 - General computation DAGs. Unlike MapReduce, which forces jobs into two rigid phases (map and reduce), Spark supports general DAGs for more flexible computation flows.
 - High-level APIs (DataFrame/Dataset). DataFrames and Datasets offer high-level abstractions that simplify working with structured data and enable query optimization.



- Unified Platform. Multiple workloads into a single platform:
 - Batch processing (similar to MapReduce)
 - Streaming (real-time data)
 - Machine learning (MLlib)
 - Graph processing (GraphX)
 - SQL queries (Spark SQL)

Spark Functionality: Core components



Resilient Distributed Datasets (RDDs):

• Distributed collections of objects (foundation for fault tolerance and parallelism.)

DataFrames and Datasets:

• Higher-level abstractions for structured and semi-structured data (Optimized via Spark's Catalyst engine).

Spark SQL:

• Query structured data using SQL.

MLlib:

• Machine learning library for scalable algorithms.

GraphX:

• Graph processing library.

Spark Functionality: Architecture



Driver Program:

• Defines the **application** and **coordinates tasks**.

Cluster Manager:

• Allocates resources (YARN, Mesos, Kubernetes).

Executors:

- Workers that execute tasks and store data partitions. **DAGs:**
- Spark builds a logical execution plan before running tasks.

Spark Functionality: Workflow



- Create RDD/DataFrame: Load data into Spark from HDFS, S3, or other sources.
- **Transformations:** Apply operations (e.g., map, filter, groupBy).
- Actions: Trigger execution (e.g., collect, save).
- **Execution:** (a) Splits tasks across nodes, (b) Uses DAG to optimize execution.

Spark: Hands on Lab

Servers Log



- Use the **COLAB** to simulate Spark basic operations
- Let's take a look into **Databricks…**

Data-Parallel DataFrame Operations

Origins of DataFrames

Recap: Data Preparation Problem

- 80% Argument: 80-90% time for finding, integrating, cleaning data
- Data scientists prefer scripting languages and in-memory libraries

Python DataFrames:

- Python pandas DataFrame for seamless data manipulations (most popular packages/features)
- DataFrame: table with a schema
- Descriptive stats and basic math, reorganization, joins, grouping, windowing
- Limitation: Only in-memory, single-node operations

import pandas as pd df = pd.read_csv('data/tmp1.csv', index_col=2) df.head() df = pd.concat(df, df[['A','C']], axis=0)

Spark DataFrames and DataSets

Overview Spark DataFrame

- DataFrame is distributed collection of rows with named/typed columns
- Relational operations (e.g., projection, selection, joins, grouping, aggregation)
- DataSources (e.g., json, jdbc, parquet, hdfs, s3, avro, hbase, csv, cassandra)

DataFrame and Dataset APIs

- DataFrame was introduced as basis for Spark SQL
- DataSets allow more customization and compile-time analysis errors (Spark 2)

DataFrame and Dataset APIs

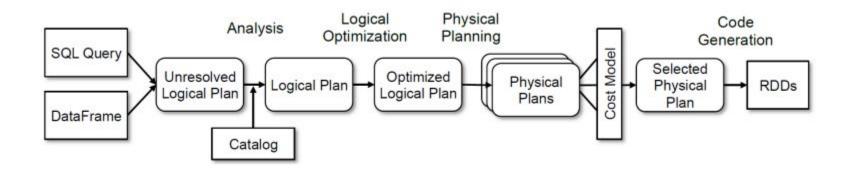
```
logs = spark.read.format("json").open("s3://logs")
logs.groupBy(logs.user_id).agg(sum(logs.time))
.write.format("jdbc").save("jdbc:mysql//...")
```

SparkSQL and DataFrame/Dataset

Overview SparkSQL

- Shark (~2013): academic prototype for SQL on Spark
- SparkSQL (~2015): reimplementation from scratch
- Common IR and compilation of SQL and DataFrame operations

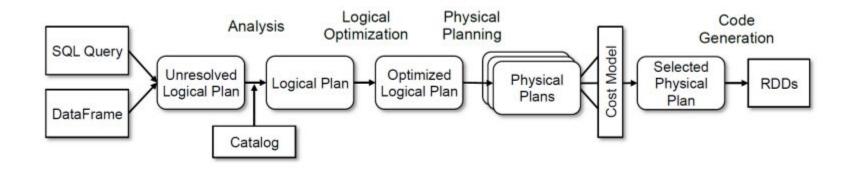
Catalyst: Query Planning



SparkSQL and DataFrame/Dataset

Performance features

- 1. Whole-stage code generation via Janino
- 2. Off-heap memory (sun.misc.Unsafe) for caching and certain operations
- 3. Pushdown of selection, projection, joins into data sources (+ join ordering)





Overview Dask

- Multi-threaded and distributed operations for arrays, bags, and dataframes
- dask.array: list of numpy n-dim arrays
- dask.dataframe: list of pandas data frames
- dask.bag:unordered list of tuples (second order functions)
- Local and distributed schedulers: threads, processes, YARN, Kubernetes, containers, HPC, and cloud, GPUs

Execution

- Lazy evaluation
- Limitation: requires static size inference
- Triggered via compute()

import dask.array as da x = da.random.random((10000,10000),
chunks=(1000,1000))
y = x + x.T
y.persist() # cache in memory
z = y[::2, 5000:].mean(axis=1) #colMeans ret
= z.compute() # returns NumPy array

Summary and Q&A

Summary and Q&A

• Summary and Q&A

- \circ $\;$ Motivation and Terminology
- \circ $\$ Data-Parallel Collection Processing
- \circ Data-Parallel DataFrame Operations

• Next Lectures

• Distributed Stream Processing [Jan 10]

Vielen Dank!